

High Performance Low Latency 16×16 bit Booth Multiplier using Novel 4-2 Compressor Structure

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ABSTRACT:

In this article, the design procedure of a low latency Booth multiplier has been proposed. With the help of a novel 4-2 compressor, a high-performance 16×16 bit Booth multiplier has been implemented, which depicts high operating frequency. To achieve this, the proposed 4-2 compressor has been utilized successively in the Partial Product Reduction Tree (PPRT) of the multiplier and by means of radix-4 Booth scheme, the multiplier has been designed. The Partial Product (PP) generation circuitry is also based on the other work published by the authors which enables the designed structure to work at the frequency of 350MHz. The main advantage of the designed compressor is the elimination of the middle stage inverters between cascaded blocks of PPRT which considerably enhances the speed of whole system. Simulation results for TSMC 0.18μm CMOS technology and 1.8V power supply have been demonstrated to confirm the correct operation of proposed 4-2 compressor. According to the results, the achieved delay of the critical path for hard test and high capacitive load, equal to 100fF, is 936ps while a power consumption of 255.15μW has been achieved at the operating frequency of 100MHZ.

KEYWORDS: Booth Multiplier, Modified Booth Encoding Scheme, 4-2 Compressor, Radix-4, Low Latency.

1. INTRODUCTION

A comprehensive literature review depicts that hardware implementation of binary multipliers contains a history of 50 years. On the other hand, implementation of algorithmic based multipliers has been launched in 1970s [1] which was right after the presentation of Booth [2], Wallace [3], and Dadda [4] multiplication algorithms. In modern terminology, this classification of multipliers which are based on a specific multiplication procedure is recognized as parallel multipliers.

Parallel multipliers and their successors, Booth multipliers, are extremely utilized in high performance communication systems such as Digital Signal Processors (DSPs) and microprocessors. In spite of high throughputs, the critical path delay in most of such systems belongs to the multiplier. Hence, the speed performance of whole structure is directly determined by the parallel multiplier used in the architecture of the system [5].

As shown in Fig. 1, the following steps can describe the multiplication process in a parallel multiplier:

- Generation of Partial Products (PPs) at the first step.

- Reduction of the products to two rows by summation in the second stage.
- Utilizing an adder structure to sum up the two remaining rows in the last stage.

The common procedure for realization of the first stage is Booth algorithm due to its higher performance in PP generation. Over the past decay, its optimized version known as Modified Booth Encoding (MBE) scheme is extremely used in high speed parallel multiplier design criteria. The main reason behind this is that with the help of MBE algorithm, the PPs will be halved while the utilized hardware is already the same as before. There are variety of structures reported in previous works for hardware realization of this method where each scheme exhibits its own benefits and drawbacks [6-13].

For implementation of the second stage, a Partial Product Reduction Tree (PPRT) is usually employed. The main tasks of this stage are summation of the PPs to reduce them to two rows. In most of the recent works, this objective has been accomplished with the help of 4-2 compressors. Considering speed performance and power dissipation as the main targets of optimization, many structures have been introduced

for hardware realization of 4-2 and higher order compressors where the emphasis was on behavioral improvement of 4-2 compressors [14-19].

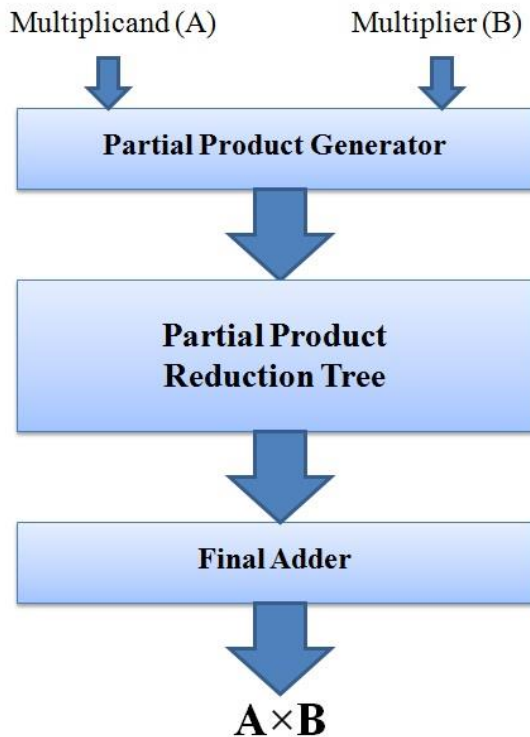


Fig. 1. Conventional scheme of a parallel multiplier.

For speed enhancement, Pass-Transistor Logic (PTL) has been employed in [14]. However, the gate level delay from inputs to the outputs was not reduced to less than 4 XOR logic gates. On the other hand, optimization of power was the main objective in [15]. Although these two circuits have shown better speed performance compared to the previous works, in 2012, two novel structures were reported in [16] and [17] which surpassed the other works from the viewpoint of gate level latency. Those architectures could achieve the gate level delay about 2 XOR logic gates by applying some modifications to the general truth table of 4-2 compressor. The recent works reported in [18] and [19], again worked on the optimization of power consumption considering new aspects such as approximate compressors.

At the final stage which pertains to the summation of two remaining rows, there are efficient adding configurations such as Carry-Lookahead Adder (CLA), Carry Select Adder (CSA), Ripple Carry Adder (RCA), and etc which can perform the task successively to finalize the multiplication process.

Taking into account that the speed improvement for 4-2 compressor is the main objective in this paper, a new scheme has been proposed which can achieve the

gate level delay of 2.25 XOR logic gates. The main idea behind the implementation of this circuit is its compatibility with the 4-2 compressor designed in [16] which enables us to construct a high performance PPRT. In order to evaluate the correct operation of the proposed PPRT, a 16×16 bit booth multiplier has been implemented. It must be mentioned that all of the circuits have been designed employing TSMC 0.18μm CMOS technology.

The paper is structured in following sections. Implementation of the proposed 4-2 compressor has been explained in section 2 while the important design considerations have been clarified for the compatibility between cascading of compressors. Section 3 illustrates the design of the 16×16 bit booth multiplier consisting of MBE scheme, PPRT and CSA. The simulation results and the comparison with previous works have been provided in section 4. Eventually, the conclusions are discussed in section 5.

2. THE PROPOSED 4-2 COMPRESSOR

As the literature states, the 4-2 compressor was firstly developed by Weinberger in 1981 and its architecture is composed of carry save adders [20]. By defining the scheme of Fig. 2 as the general architecture, to design a 4-2 compressor, two Full Adders (FAs) should be cascaded. As a result, the main difference between a 4-2 compressor and an FA can be distinguished by the horizontal path (C_{out}) [18]. The two vertical paths (Sum and $Carry$) will be the common point between these two architectures. Table 1 describes the corresponding truth table for hardware level implementation of a 4-2 compressor.

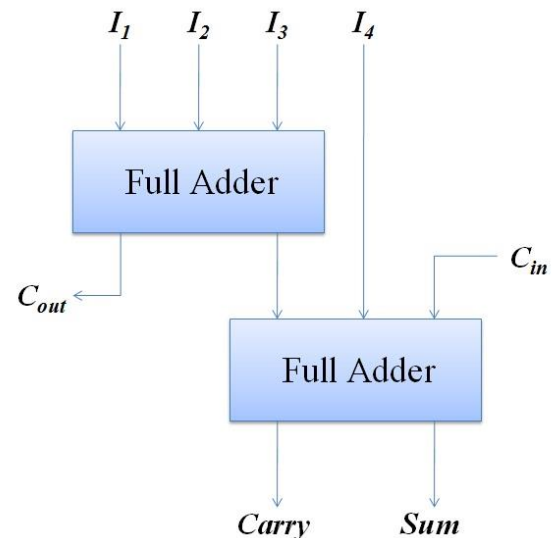


Fig. 2. Conventional architecture of 4-2 compressor.

Table 1. General truth table of 4-2 compressor.

n	C_{in}	C_{out}	Carry	Sum
0	0	0	0	0
1	0	0	0	1
2	0	-	-	0
3	0	1	0	1
4	0	1	1	0
0	1	0	0	1
1	1	0	1	0
2	1	-	-	1
3	1	1	1	0
4	1	1	1	1

Investigation of conventional truth table depicts that there are two neutral states for C_{out} and Carry outputs and these states exactly happen when two of four input bits (I_1, I_2, I_3 and I_4) in Fig. 2 have logical value of one. The authors of [16] and [17] have assumed some conditions for these two states which led to the simplification of this table. The result of this effort was two high speed circuits which show gate level delay equal to almost 2 XOR logic gates from inputs to the outputs while the structure reported in [16] which is shown in Fig. 3 has better speed performance.

However, in Fig. 3, all of the outputs have been obtained in complementary state. Therefore, an inverter gate after each output is needed so that the real state of the outputs can be achieved. Therefore, the real gate level delay will increase to 2.5 XOR logic gates.

To solve this problem, the need for a compatible structure becomes necessary so that the corresponding inputs of such structure can operate in complementary state. This was the motif for implementation of a novel 4-2 compressor in which its inputs are driven in complementary mode. With the help of such configuration, there is no need for any inverters at output nodes of Fig. 3 which maintains the high-speed performance of the circuit.

The proposed architecture has been shown in Fig. 4 where all of the inputs are in complementary state. By defining the control signals E and F as:

$$E = (I_1 \oplus I_2) \oplus (I_3 \oplus I_4) \tag{1}$$

$$F = \overline{I_1 + I_2 + I_3 + I_4} \tag{2}$$

the Carry output will abide by:

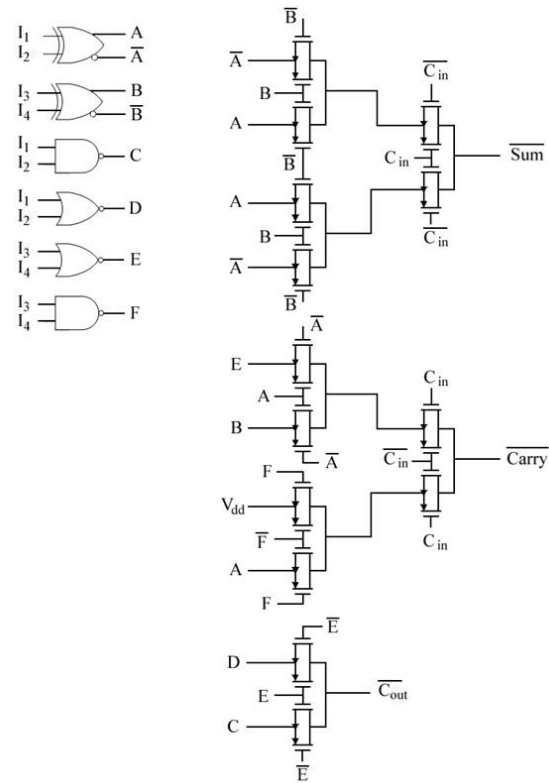


Fig. 3. 4-2 compressor designed in [16].

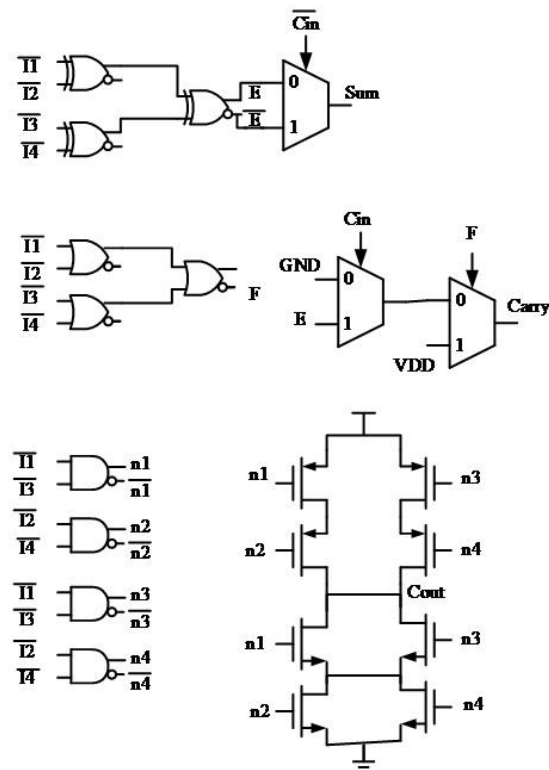


Fig. 4. The proposed 4-2 compressor.

$$Carry = (E \cdot C_{in} \cdot \bar{F}) + F \tag{3}$$

Furthermore, based on Table 1, if at least 2 of inputs are equal to logic “1” value, then C_{out} will rise to high logic level. Finally, for the *Sum* output we have:

$$Sum = (I_1 \oplus I_2) \oplus (I_3 \oplus I_4) \oplus C_{in} \tag{4}$$

To calculate the gate level delay for proposed circuit, as Fig. 4 demonstrates, a latency of 2 XOR logic gates plus one transistor is expectable. Furthermore, by defining the latency of channel-ready transistor equal to 0.25, then the latency can be considered as 2.25 XOR logic gates.

3. THE MULTIPLIER

To show the efficiency of proposed 4-2 compressor in conjunction with its compatibility with 4-2 compressor of Fig. 3, a 16×16 bit Booth multiplier utilizing the MBE scheme of [22] (which was previously designed and published by the authors), has been implemented.

Fig. 5 shows the proposed concept for implementation of such multiplier in which an MBE scheme along with two cascaded 4-2 compressor structures should be utilized to generate two rows of bits for final addition. The final adder will sum these rows so that the final product can be obtained.



Fig. 5. Proposed concept for implementation of 16×16 bit Booth multiplier.

To achieve this, the circuitry of MBE was redesigned in CMOS 0.18μm technology considering the truth table defined in Table 2.

According to this truth table, the parameters *X*, *Z* and *Neg* are utilized for radix-4 bit encoding. For simplification, \bar{Z} is generated at the encoder section due to its simpler Boolean expression [22].

Table 2. The truth table of Booth scheme.

d_i	b_{2i+1}	b_{2i}	b_{2i-1}	<i>Neg</i>	<i>X</i>	<i>Z</i>
0	0	0	0	0	0	0
1	0	0	1	0	1	1
1	0	1	0	0	1	1
2	0	1	1	0	0	1
-2	1	0	0	1	0	1
-1	1	0	1	1	1	1
-1	1	1	0	1	1	1
-0	1	1	1	1	0	0

Fig. 6 shows the encoder section while in Fig. 7, the decoder block has been demonstrated for generation of PP rows.

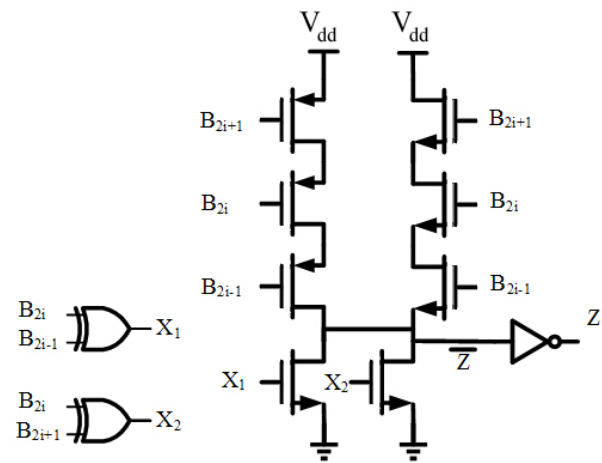


Fig. 6. Redesigned circuit of Booth encoder from [22].

The main advantage of this scheme is its enhanced speed for PP generation compared with architecture of [10]. A critical path delay less than 2 XOR logic gates is expected for the employed MBE scheme. In addition, because of the uniform paths, there would not be any side effects of glitch at the output waveforms.

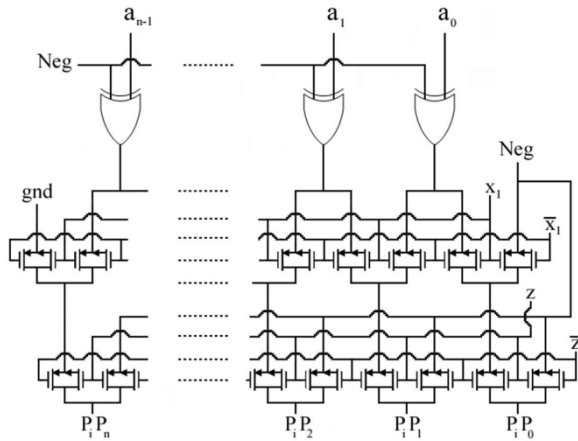


Fig. 7. Redesigned circuit of Booth decoder from [22].

Following the concept of achieving higher operating frequency, CSA has been employed at the final addition stage to sum the final remaining rows [23]. Fig. 8 shows the basic structure used for implementation of CSA architecture which utilizes four bit adders.

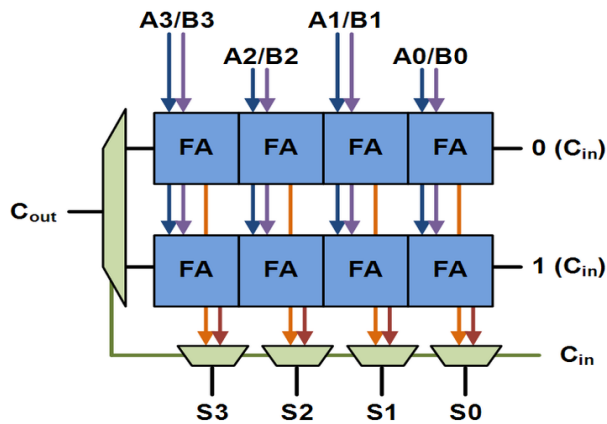


Fig. 8. Employed CSA for implementation of final adder.

As shown in Fig. 8, two 4-bit adder blocks are employed in the structure of CSA. One of the adders presumes that the input carry value is zero, while the other considers the logic value of one for the input carry. The carry outputs of these blocks are multiplexed together to produce the final carry output of the architecture. The multiplexing is performed by selecting which adder contains the accurate value.

In addition, Fig. 9 illustrates the architecture of FA utilized in the configuration of CSA in which Fig. 9(a) presents the proposed circuit and Fig. 9(b) shows the modified structure. To avoid cascading of the multiplexers in CSA arrays which degrades the speed performance, for odd FA cells, the proposed circuit of Fig. 9(a) has been used while for even cells the modified structure of Fig. 9(b) was employed.

Considering that the first eight bits along with the last four bits of multiplication process in the implemented 16×16 bit multiplier can be created at the PPRT, the proposed architecture of CSA for summation of twenty bits has been shown in Fig. 10.

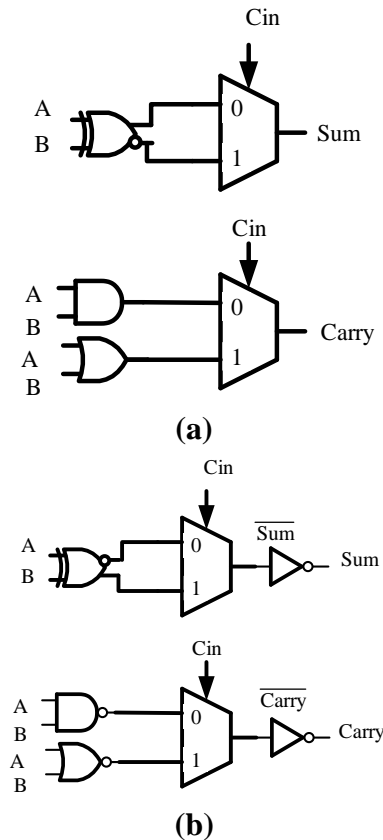


Fig. 9. Proposed FA scheme (a) main circuit (b) modified circuit.

Finally, Fig. 11 illustrates the scheme of multiplication in which for the middle stages, the cascaded 4-2 compressors are the structures shown in Fig. 3 and Fig. 4, respectively. As described above, the first eight bits and the last four bits can be computed in the PPRT.

Therefore, the delay of critical path can be calculated by means of the following relation:

$$T_{delay} = t_{Booth} + t_{Comp1} + t_{Comp2} + t_{Adder} \quad (5)$$

In which, t_{Booth} describes the latency for generation of PPs by MBE structure. In addition, t_{Comp1} defines the propagation delay for 4-2 compressor shown in Fig. 3, while t_{Comp2} represents the latency associated with 4-2 compressor of Fig. 4. Finally, t_{Adder} denotes the delay of final addition which is implemented by means of FAs.

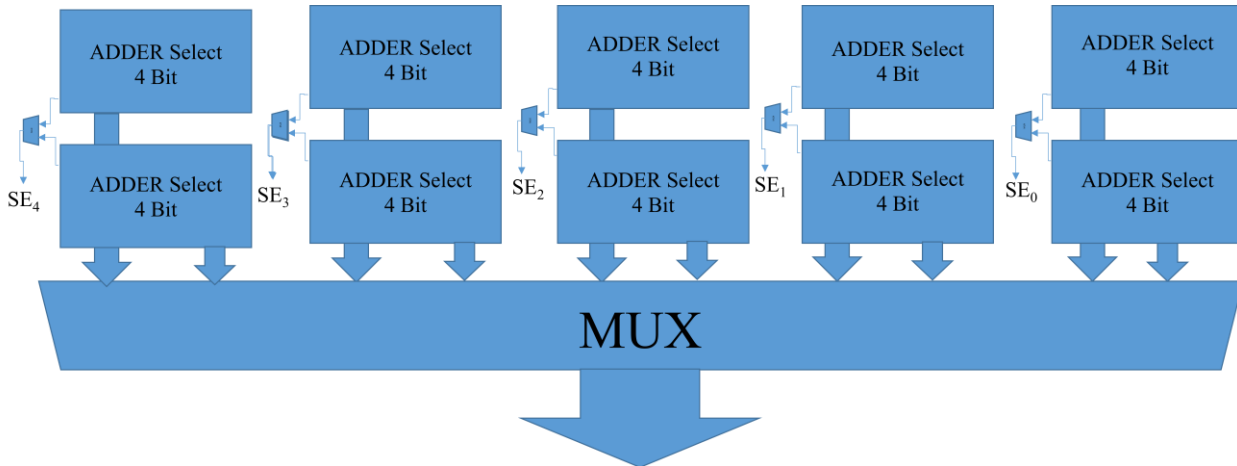


Fig. 10. The scheme of implemented CSA for 20 bit addition.

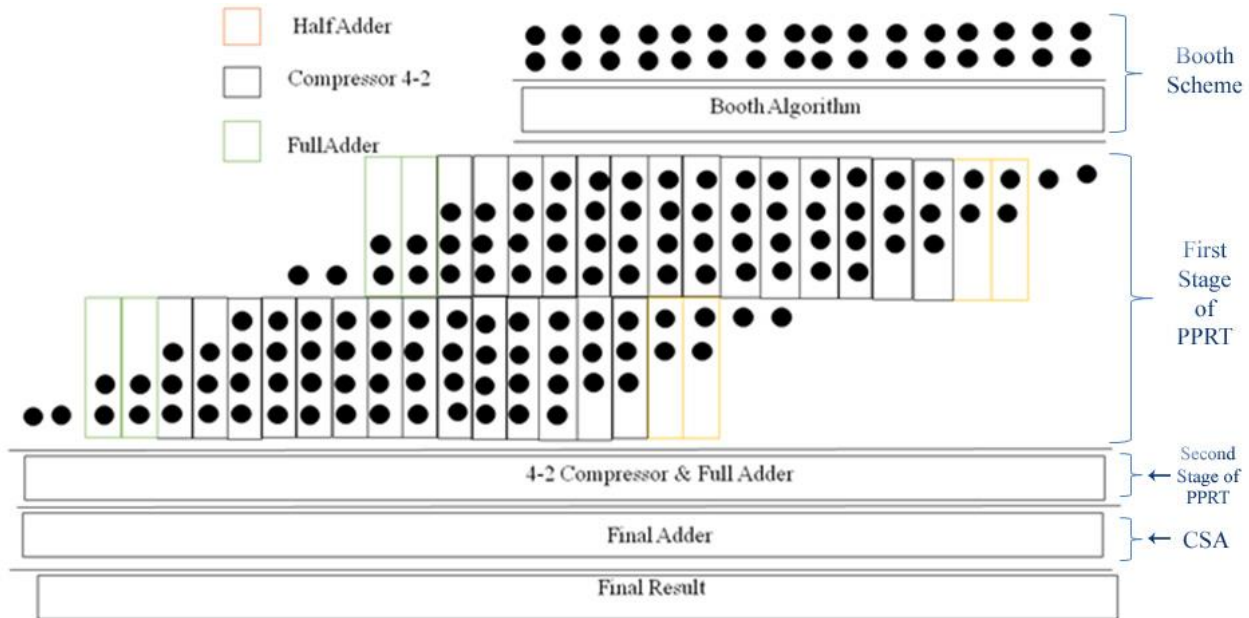


Fig. 11. The scheme of proposed 16x16 bit Booth multiplier.

It must be mentioned that in (5), the loading of each block will directly affect the latency of its previous level and it has been considered for accurate measurement of propagation delay in the proposed multiplier.

4. SIMULATION RESULTS AND COMPARISON

In order to investigate the correct operation of proposed 4-2 compressor and obtain the delay of the critical path, simulations using HSPICE for TSMC standard 0.18µm CMOS technology and 1.8V power

supply were performed in which all of the outputs were followed with high capacitive loads equal to 100fF. Moreover, to perform a hard test for the designed circuitry, the arrangement of the input bits was set in a way so that the critical path will surely be enabled.

Fig. 12 depicts the result which confirms the accurate operation of implemented structure. It must be mentioned that the obtained delay which is equal to 936ps has been obtained for the high capacitive load mentioned before. Moreover, the power consumption at the operating frequency of 100MHz was 255.15 µW.

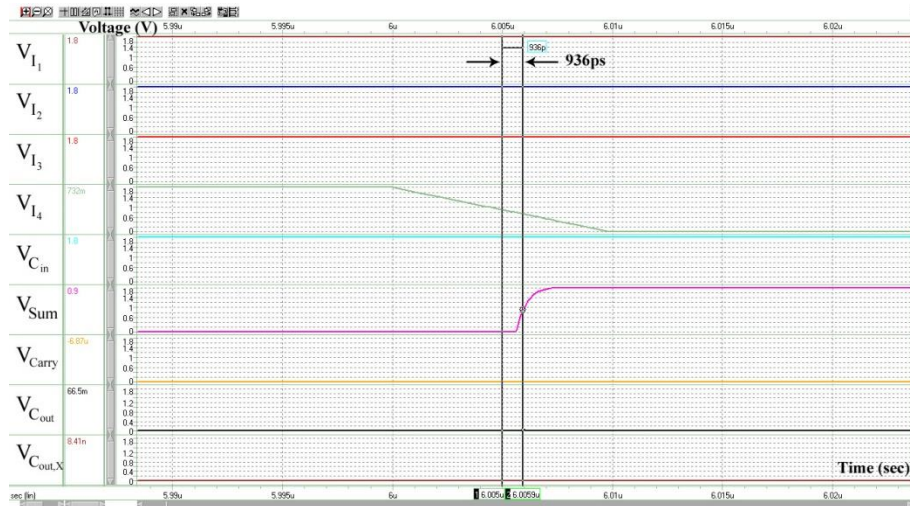


Fig. 12. Simulation results for delay calculation of proposed compressor.

For a fair comparison between proposed 4-2 compressor and the previously reported works, the compressor architectures of [14], [15], [16] and [17] were also redesigned employing the gates used for implementation of our circuitry.

These works are then simulated considering the same conditions mentioned above. Based on the results which have been presented in Fig. 13 and Fig. 14, the delay of proposed compressor is lower than the previous works while the power consumption is favorably comparable with those works.

Considering the obtained results from simulations, Table 3 illustrates the comparison of the proposed work with previous designs. It must be mentioned that the results for [16] has been obtained considering one inverter cascading at the output nodes. Hence, the lowest latency will belong to the proposed structure.

Finally, to obtain the delay of the critical path for the proposed 16×16 bit multiplier, the layout of the proposed multiplier was drawn by means of TSMC 0.18μm CMOS technology file where it is shown in Fig. 15.

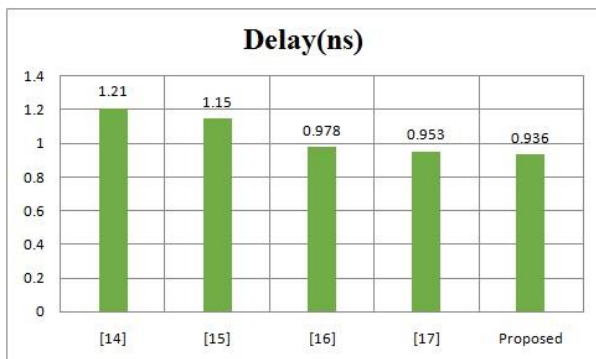


Fig. 13. Delay comparison of proposed work and previous designs.

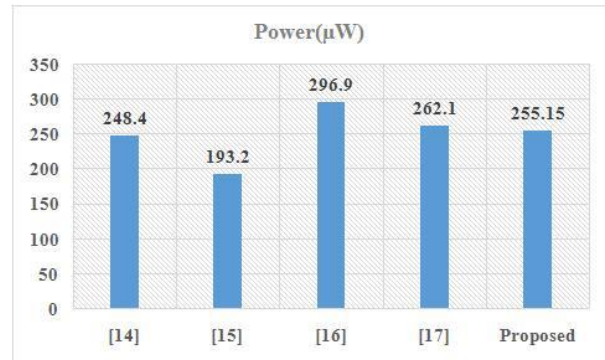


Fig. 14. The comparison of power dissipation between proposed work and previous designs.

Table 3. Comparison of the proposed work and previous designs

Work	Proposed	[14]	[15]	[16]	[17]
Technology (μm)	0.18	0.18	0.18	0.18	0.18
Transistor Count	68	64	60	74	70
Gate Level Delay (XOR)	2.25 gates	4 gates	3 gates	2.5 gates	2.25 gates
Power (μW) @100MHz	255.15	248.4	193.2	296.9	262.1
Delay(ps)	936	1210	1150	978	953

An active area equal to 197μm× 289μm is consumed for the whole architecture. After that, the parasitics were extracted, post-layout simulations using HSPICE for TSMC standard 0.18μm CMOS technology and 1.8V power supply were carried out for the designed multiplier architecture.

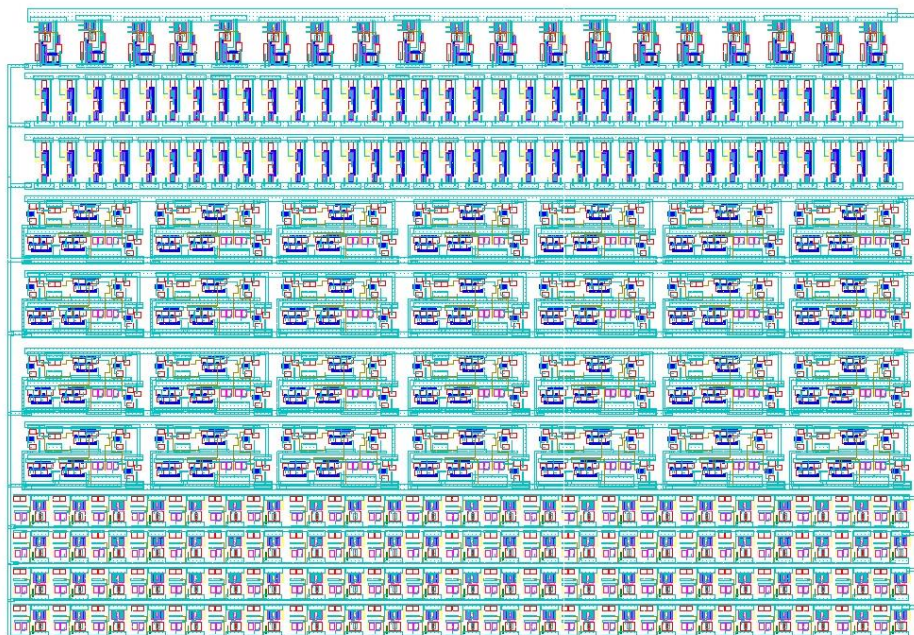


Fig.15. Layout of the proposed 16×16 bit Booth multiplier.

Table 4 summarizes the obtained results for propagation delay of implemented multiplier along with the corresponding latencies of different blocks utilized for realization of the multiplier. Based on the results, the total latency of the proposed multiplier is equal to 2841ps which confirms that the implemented architecture can operate at the frequencies up to 351MHz if no pipeline stage is utilized for the speed enhancement.

Table 4. Propagation delay based on post-layout simulation results for the proposed multiplier and its building blocks.

MBE Structure	255ps
4-2 compressor of Fig. 3	237ps
4-2 compressor of Fig. 4	249ps
Final Adder	2100ps
Total Latency	2841ps

Meanwhile, according to the results, the summation of four latencies obtained for MBE structure, 4-2 compressors, and the final adder, will exactly become equal to the obtained value for the total latency of the proposed multiplier. This also complies with the obtained value in (5).

It must be mentioned that for the pipeline implementation of this structure much higher operating frequencies are expected, since the clocked feature of

pipeline architecture will enable the subsections to operate up to their own frequency limits. For instance, the proposed 4-2 compressor can work in the frequency of 1GHz even for high capacitive loads.

Moreover, the differences between the obtained delays of the compressors in Table 3 and Table 4 are due to the different capacitive loads. For the results of Table 3, high capacitive load is used for hard test which significantly increases the latency.

5. CONCLUSION

A novel architecture for 4-2 compressor has been proposed in this paper which demonstrates excellent speed performance. Having the compatibility feature with 4-2 compressor circuit reported in [16], the designed structure can successively be cascaded with that circuit to construct a high speed PPRT for utilization in high performance parallel multipliers. Considering the fact that total transistor count of the proposed 4-2 compressor is almost the same as previous works, its active area consumption will have no significant difference compared with those structures.

Based on the simulation results for TSMC 0.18μm CMOS technology and 1.8V power supply, the obtained delays of the proposed 4-2 compressor for normal and high capacitive loads are 249ps and 936ps, respectively.

In addition, the designed Booth multiplier which is based on PPRT consisting of combined 4-2 compressors can operate at the frequencies up to 351MHz.

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